INSTITUTE OF AERONAUTICAL ENGINEERING

(Autonomous) Dundigal-500043, Hyderabad

B.Tech IV SEMESTER END EXAMINATIONS (REGULAR / SUPPLEMENTARY) - AUGUST 2023

Regulation: UG-20

Time: 3 Hours

THEORY OF COMPUTATION (CSE | CSE(AI&ML) | CSIT | IT)

Max Marks: 70

Answer ALL questions in Module I and II Answer ONE out of two questions in Modules III, IV and V All Questions Carry Equal Marks All parts of the question must be answered in one place only

$\mathbf{MODULE}-\mathbf{I}$

- 1. (a) Describe how various rules used for construction of ϵ -NFA with suitable example and also write th steps for NFA with ϵ to NFA conversion with an example. [BL: Understand| CO: 1|Marks: 7]
 - (b) Differentiate DFA and NFA with an example. Give DFA's accepting the following languages over the alphabet $\{0, 1\}$.

i) The set of all strings ending in 00.

 $S \rightarrow aS/bA/b$, $A \rightarrow aA/bS/a$

ii) The set of all strings with three consecutive 0's. (not necessarily at the end).

[BL: Apply] CO: 1|Marks: 7]

$\mathbf{MODULE}-\mathbf{II}$

2. (a) Illustrate the steps for conversion of regular grammar to finite automata? Construct the FA for the following grammar

[BL: Understand] CO: 2|Marks: 7]

(b) Convert the following NFA given in Figure 1 into a RE using the state elimination method.

[BL: Apply| CO: 2|Marks: 7]

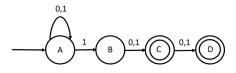


Figure 1

$\mathbf{MODULE}-\mathbf{III}$

3. (a) Enlist any three closure properties of regular languages with an illustration.

[BL: Understand| CO: 3|Marks: 7]

- (b) Prove that the given languages are not regular using the pumping lemma.
 i) L = {w = a^Jb^K : J = K}
 ii) L = {w = a^Jb^K : J > K} [BL: Apply| CO: 3|Marks: 7]
- 4. (a) Summarize about ambiguous Grammar. Check whether the grammar: S \rightarrow aAB, A \rightarrow bC/cd, C \rightarrow cd, B \rightarrow c/d is ambiguous or not? [BL: Understand| CO: 4|Marks: 7]

(b) Convert the CFG into CNF S \rightarrow AB | Aa A \rightarrow aaA | a B \rightarrow bbB | b

[BL: Apply] CO: 4|Marks: 7]

MODULE - IV

- 5. (a) Differentiate push down automata (PDA) and DPDA with an example. Show that for every PDA there exists a CFG such that L(G) = N(P). [BL: Understand CO: 5|Marks: 7]
 - (b) Construct a PDA that accepts the language $L = \{a^n b^m c^m d^n | m, n \ge 1\}$ by using an empty stack. [BL: Apply] CO: 5|Marks: 7]
- 6. (a) Outline the followings:
 - i) Closure properties of context free languages.
 - ii) Deterministic context free languages
 - iii) Deterministic Pushdown Automata.
 - (b) Convert the grammar CFG to a PDA $E \rightarrow E + E$

 $E \rightarrow id$

[BL: Apply] CO: 5|Marks: 7]

[BL: Understand] CO: 5 Marks: 7]

MODULE - V

- 7. (a) Give an overview of recursively enumerable language. Mention the properties of recursive and non-recursive enumerable languages. [BL: Understand| CO: 6|Marks: 7]
 - (b) Construct a Turing machine that will accept the language consists of all palindromes of 0's and 1's? [BL: Apply| CO: 6|Marks: 7]
- 8. (a) Describe in detail about different types of Turing machines Compare and contrast the finite state machine, PDA and Turing machine. [BL: Understand| CO: 6|Marks: 7]

(b) Build Turing machine M, 1011, where $M = (q_1, q_2, q_3, 0, 1, 0, 1, B, \delta, q_1, B, q_2)$ has moves, $\delta(q_1, 1) = (q_3, 0, R)$ $\delta(q_3, 0) = (q_1, 1, R)$ $\delta(q_3, 1) = (q_2, 0, R)$ $\delta(q_3, B) = (q_3, 1, L)$

[BL: Apply] CO: 6|Marks: 7]

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