## INSTITUTE OF AERONAUTICAL ENGINEERING (Autonomous)

Dundigal, Hyderabad -500 043

## INFORMATION TECHNOLOGY <br> TUTORIAL QUESTION BANK

| Course Title | THEORY OF COMPUTATION |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Course Code | AITB03 |  |  |  |  |
| Programme | B. Tech |  |  |  |  |
| Semester | IV IT |  |  |  |  |
| Year | 2019-2020 |  |  |  |  |
| Course type | Core |  |  |  |  |
| Regulation | IARE - R18 |  |  |  |  |
| Course Structure | Theory |  |  | Practical |  |
|  | Lectures | Tutorials | Credits | Laboratory | Credits |
|  | 3 | 1 | 4 | - | - |
| Chief coordinate | Dr.K Srinivasa Reddy, Associate Professor |  |  |  |  |
| Course Faculty | Dr.K Srinivasa Reddy, Associate Professor |  |  |  |  |

## COURSE OVERVIEW:

In theoretical computer science and mathematics, the theory of computation is the branch that deals with how efficiently problems can be solved on a model of computation, using an algorithm. The field is divided into three major branches automata theory and languages, computability theory, and computational complexity theory, which are linked by the question: "What are the fundamental capabilities and limitations of computers?"

## COURSE OBJECTIVES:

## The course should enable the students to:

| I. | Comprehend abstract, mathematical models of computation and use them to solve computational <br> problems |
| :--- | :--- |
| II. | Interpret the relationship between formal languages in Chomsky's hierarchy and different machines. |
| III. | Analyze and explain the behavior of push-down automata. |
| IV. | Understand the limits and capacities of Turing's machines to recognize languages. |

## COURSE OUTCOMES (COs):

| COs | Course Outcome |
| :--- | :--- |
| CO 1 | Understand the functionality of deterministic finite automata and Non-deterministic finite automata |
| $\mathbf{C O}$ 2 | Apply the regular languages , regular expressions to construct finite automata |
| $\mathbf{C O} 3$ | Apply the context free grammars to construct derivation trees and the accept various strings |
| $\mathbf{C O} 4$ | Compare the functionality of push down automata with deterministic finite automata |
| $\mathbf{C O} 5$ | Apply the concept of Turing machines to solve the complex functions |

## COURSE LEARNING OUTCOMES:

Students, who complete the course, will have demonstrated the asking to do the following:

| AITB03.01 | Use the definitions and notations for sets, relations and functions in defining and study Finite <br> Automata |
| :--- | :--- |
| AITB03.02 | Remember on formal languages and Kleene"s Theorem to intend programming languages |
| AITB03.03 | Construct deterministic and nondeterministic finite state automata (DFA and NFA) for solving <br> simple decision problems. |
| AITB03.04 | Perform conversions between nondeterministic finite automata, deterministic finite automata and regular <br> expressions and finite state automata to gain Remember about formal proofs in computer science |
| AITB03.05 | Remember on recursive definitions of regular languages, regular expressions and the use of regular <br> expressions to represent regular languages |
| AITB03.06 | Detailed Remember on the relationship between regular expressions and finite automata |
| AITB03.07 | Identify that few languages are not regular by using Pumping lemma |
| AITB03.08 | Remember on Left Linear grammar, Right Linear grammars and converting grammars into Finite <br> Automata. |
| AITB03.09 | Understand the fundamental role played by Context-Free Grammars (CFG) in designing formal <br> computer languages with simple examples |
| AITB03.10 | Remember on Context Free Grammars so that able to prove properties of Context Free Grammars. |
| AITB03.11 | Identify relationship between regular languages and context-free grammars |
| AITB03.12 | Use the pumping lemma for Context Free Languages to show that a language is not context-free |
| AITB03.13 | Understand the equivalence between Context-Free Grammars and Non-deterministic Pushdown <br> Automata |
| AITB03.14 | Understand deterministic Pushdown Automata to parse formal language strings by using (i) top down or <br> (ii) bottom up techniques |
| AITB03.15 | Remember on converting Context-Free Grammars into pushdown automata to identify the acceptance <br> of a string by the Context Free Language |
| AITB03.16 | Understand the path processing computation using Turing Machines (Deterministic and Non- <br> Deterministic) and Church-Turing Thesis in computers. |
| AITB03.17 | Remember on non-halting Turing Machine accepted by Recursively Enumerable Languages |
| AITB03.18 | Understand the power of the Turing Machine, as an abstract automaton, that describes <br> computation, effectively and efficiently |
| AITB03.20 | Theory of Computation is important in programming language design, parsers, web-Scrappers, Natural <br> Language Processing (NLP), and is at the heart of modern compiler architectures. <br> competitive exams. |

## TUTORIAL QUESTION BANK

| MODULE-1 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| FINITE AUTOMATA |  |  |  |  |
| PART - A (Short Answer Questions) |  |  |  |  |
| S. No | Questions | $\begin{array}{c\|} \hline \text { Blooms } \\ \text { Taxonomy } \\ \text { Level } \\ \hline \end{array}$ | Course Outcomes | Course <br> Learning <br> Outcomes |
| 1 | Demonstrate finite state Automata. | Understand | CO1 | AITB03.01 |
| 2 | Distinguish between DFA and NFA. | analyze | CO1 | AITB03.03 |
| 3 | Describe the concept of String. | Understand | CO1 | AITB03.01 |
| 4 | Describe transition function of DFA. | Understand | CO1 | AITB03.01 |
| 5 | Demonstrate about $\varepsilon$-transitions. | Understand | CO1 | AITB03.03 |
| 6 | Highlight the power of an alphabet ( $\Sigma^{*}$ ). | Understand | CO1 | AITB03.01 |
| 7 | List out the applications of finite automata. | Remember | CO1 | AITB03.01 |
| 8 | Describe the concept of Null string. | Understand | CO1 | AITB03.01 |
| 9 | Demonstrate the Kleene Star? | Remember | CO1 | AITB03.01 |
| 10 | Illustrate NFA with example. | Understand | CO1 | AITB03.03 |
| 11 | Describe transition diagram for DFA accepting string ending with 00 | Understand | CO1 | AITB03.01 |
| 12 | Construct DFA for a string accepting odd number of 0"s. | apply | CO1 | AITB03.03 |
| 13 | Illustrate transition diagram for DFA to accept exactly one ,„" defined. over an alphabet $\sum=\{\mathrm{a}, \mathrm{b}\}$.. | Understand | CO1 | AITB03.03 |
| 14 | Construct DFA for odd number of 1"s. | Apply | CO1 | AITB03.03 |
| 15 | Demonstrate $\varepsilon$ - closure. | Understand | CO1 | AITB03.03 |
| 16 | Describe FSM and its structure with an example. | Understand | CO1 | AITB03.01 |
| 17 | State the Mathematical definition of Finite Automata. | Remember | CO1 | AITB03.01 |
| 18 | Construct DFA for even number of 1"s. | Apply | CO1 | AITB03.01 |
| 19 | Demonstrate DFA mathematically. | Remember | CO1 | AITB03.01 |
| 20 | Construct DFA for the language accepting strings which contains 001 as substring. | Apply | CO1 | AITB03.03 |
| Part - B (Long Answer Questions) |  |  |  |  |
| 1 | Construct a DFA to accept set of all strings ending with 0101. | Apply | CO1 | AITB03.01 |
| 2 | Evaluate the DFA with the set of strings having ,,aaa" as a substring over an alphabet $\sum=\{\mathrm{a}, \mathrm{b}\}$. | Evaluate | CO1 | AITB03.01 |
| 3 | List out the various differences between DFA and NFA | Remember | CO1 | AITB03.01 |
| 4 | Describe NFA with $\varepsilon$ to NFA conversion with an example. | Understand | CO1 | AITB03.01 |
| 5 | Construct a DFA to accept the string a"s and b"s ending with abb over an alphabet $\sum=\{\mathrm{a}, \mathrm{b}\}$ | Apply | CO1 | AITB03.03 |
| 6 | Describe the properties and operations of strings and languages. | Understand | CO1 | AITB03.01 |
| 7 | Construct a DFA that any given decimal number is divisible by 3 . | Apply | CO1 | AITB03.04 |
| 8 | Design DFA for the following languages shown below $\Sigma=\{\mathrm{a}, \mathrm{b}\}$ <br> a) $\mathrm{L}=\{\mathrm{w} / \mathrm{w}$ is any string that doesn"t contain exactly two a"s $\}$ <br> b) $\mathrm{L}=\left\{\mathrm{w} / \mathrm{w}\right.$ is any string that contain atmost $\left.3 \mathrm{a}^{\text {" }} \mathrm{s}\right\}$ | Apply | CO1 | AITB03.03 |
| 9 | Construct the following NFA with $\varepsilon$ to NFA. | Apply | CO1 | AITB03.03 |
| 10 | Construct Finite Automata and draw FA for the strings over an alphabet $\Sigma=\{0,1\}$ <br> (i) The string with even no of 0 "s and odd no of 1 "s | Apply | CO1 | AITB03.03 |


|  | (ii) The string with odd no of 0"s and odd no of 1"s |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 11 | Construct a DFA, the language recognized by the Automaton being $\mathrm{L}=\{\mathrm{w} / \mathrm{w}$ contains neither the substring ab nor ba $\}$.Draw the transition table. | Apply | CO1 | AITB03.04 |
| 12 | Convert the following NFA into DFA. | Understand | CO1 | AITB03.04 |
| 13 | $\begin{aligned} & \text { Construct the DFA for the following language } \\ & \mathrm{L}=\left\{\mathrm{w} / \mathrm{w} \mid \bmod 3=0, \mathrm{w} \text { belongs to }(\mathrm{a}, \mathrm{~b})^{*}\right\} \\ & \mathrm{L}=\left\{\mathrm{w} / \mathrm{w} \mid \bmod 3=1, \mathrm{w} \text { belongs to }(\mathrm{a}, \mathrm{~b})^{*}\right\} \end{aligned}$ | Apply | CO1 | AITB03.04 |
| 14 | Design a DFA for the following language over an alphabet $\sum=\{0,1\}$ <br> i) The string with even no of 0 "s and even no of 1 "s <br> ii) The string with odd no of 0 "s and even no of 1 "s | Create | CO1 | AITB03.03 |
| 15 | Convert the following NFA into equivalent DFA. | Understand | CO1 | AITB03.03 |
| 16 | Convert the following NFA- $\varepsilon$ to NFA. | Analyze | CO1 | AITB03.03 |
| 17 | ```Construct the DFA for the following language \(\mathrm{L}=\left\{\mathrm{w} / \mathrm{n}_{\mathrm{a}}\|\mathrm{w}| \bmod 3=0, \mathrm{w}\right.\) belongs to \(\left.(\mathrm{a}, \mathrm{b})^{*}\right\}\) \(\mathrm{L}=\left\{\mathrm{w} / \mathrm{n}_{\mathrm{a}}|\mathrm{w}| \bmod 3=1, \mathrm{w}\right.\) belongs to \(\left.(\mathrm{a}, \mathrm{b})^{*}\right\}\)``` | Apply | CO1 | AITB03.03 |
| 18 | Construct the following NFA with $\varepsilon$ to NFA. | Apply | CO1 | AITB03.03 |
| 19 | Construct a DFA that accepts set of strings starts with 01and ends with 01 over alphabet $\sum=\{0,1\}$ | Apply | CO1 | AITB03.03 |
| 20 | Illustrate the model and behavior of finite automata with neat block diagram. | Understand | CO1 | AITB03.04 |
|  | Part - C (Problem Solving and Critical Thinking | Questions) |  |  |
| 1 | Design NFA for accepting any binary string that contains 11 as a substring and Convert to DFA. | Create | CO1 | AITB03.04 |
| 2 | Construct NFA with $\varepsilon$ to equivalent DFA | Understand | CO1 | AITB03.04 |


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| :---: | :---: | :---: | :---: | :---: |
| 3 | Construct a DFA that any given decimal number is divisible by 5 . | Apply | CO1 | AITB03.04 |
| 4 | Construct the DFA for the following language <br> $\mathrm{L}=\left\{\mathrm{w} /\|\mathrm{w}\| \bmod 5=0, \mathrm{w}\right.$ belongs to $\left.(\mathrm{a}, \mathrm{b})^{*}\right\}$ <br> $\mathrm{L}=\left\{\mathrm{w} /\|\mathrm{w}\| \bmod 5=1, \mathrm{w}\right.$ belongs to $\left.(\mathrm{a}, \mathrm{b})^{*}\right\}$ |  |  |  |
| 5 | Design the NFA from a given NFA with $\varepsilon$ machine $\mathrm{M}=\left(\left\{\mathrm{q}_{0}, \mathrm{q}_{1}, \mathrm{q}_{2}\right\},\{0,1,2\}, \delta, \mathrm{q}_{0},\left\{\mathrm{q}_{2}\right\}\right)$ where $\delta$ is given by $\left[\delta\left(\mathrm{q}_{0}, 0\right)=\left\{\mathrm{q}_{0}\right\}, \delta\left(\mathrm{q}_{0}, 1\right)=\phi, \delta\left(\mathrm{q}_{0}, 2\right)=\phi, \delta\left(\mathrm{q}_{0}, \varepsilon\right)=\left\{\mathrm{q}_{1}\right\}\right]$ $\left[\delta\left(\mathrm{q}_{1}, 0\right)=\phi, \delta\left(\mathrm{q}_{1}, 1\right)=\left\{\mathrm{q}_{1}\right\}, \delta\left(\mathrm{q}_{1}, 2\right)=\phi, \delta\left(\mathrm{q}_{1}, \varepsilon\right)=\left\{\mathrm{q}_{2}\right\}\right]$ $\left[\delta\left(q_{2}, 0\right)=\phi, \delta\left(q_{2}, 1\right)=\phi, \delta\left(q_{2}, 2\right)=\left\{q_{2}\right\}, \delta\left(q_{2}, \varepsilon\right)=\phi\right]$ | Create | CO1 | AITB03.03 |
| 6 | Construct a NFA that strings such that the third symbol from the right end is a 0 over an alphabet $\sum=\{0,1\}$. And Convert it into equivalent DFA. | Apply | CO1 | AITB03.03 |
| 7 |  | Apply | CO1 | AITB03.04 |
| 8 | Develop the transition diagram for the below NFA and then convert its equivalent transitition diagram for DFA. | Apply | CO1 | AITB03.03 |
| 9 | Construct the DFA that will accept those words from $\sum=\{a, b\}$ where the number of a's is divisible by two and the number of b's is divisible by three. Sketch the transition table of the finite automata. | Apply | CO1 | AITB03.03 |
| 10 | Construct the DFA that will accept those words from alphabets $\Sigma=\{\mathrm{a}, \mathrm{b}\}$ where the number of b " s is divisible by three. Sketch the transition table and diagram of the finite Automata. | Apply | CO1 | AITB03.04 |
|  | MODULE-II |  |  |  |
|  | REGULAR LANGUAGES |  |  |  |
|  | PART - A (Short Answer Ques | ons) |  |  |
| S. No | Questions | Blooms Taxonomy Level | Course Outcomes | Course Learning Outcomes |
| 1 | Demonstrate Regular Languages. | Understand | CO2 | AITB03.05 |
| 2 | List out any two applications of regular expression. | Remember | CO2 | AITB03.05 |
| 3 | Construct Pumping Lemma for Regular Languages. | Apply | CO2 | AITB03.07 |
| 4 | Illustrate an example for a regular set? | Remember | CO2 | AITB03.05 |


| 5 | Construct the Regular Expression for the empty string. | Remember | CO2 | A ITB03.08 |
| :---: | :---: | :---: | :---: | :---: |
| 6 | Describe regular expression for denoting language containing empty | Understand | CO 2 | AITB03.05 |
| 7 | Demonstrate right linear grammars. | Understand | CO 2 | AITB03.08 |
| 8 | Construct the Regular Expression for the set of binary strings. | Apply | CO2 | AITB03.05 |
| 9 | Demonstrate Regular grammars. | Understand | CO2 | AITB03.05 |
| 10 | List out the advantages of regular expressions. | Remember | CO2 | AITB03.05 |
| 11 | Demonstrate Regular set? | Understand | CO2 | AITB03.05 |
| 12 | State regular expressions for the Set of strings over $\{0,1\}$ whose last two symbols are the same. | Remember | CO2 | AITB03.05 |
| 13 | Describe the regular language generated by regular expression $(0+1) * 001(0+1)^{*}$. | Understand | CO2 | AITB03.05 |
| 14 | Summarize the difference between left linear and right linear | Understand | CO2 | AITB03.08 |
| 15 | Describe the Regular Expression to generate at least one b over $\Sigma=\{\mathrm{a}, \mathrm{b}\}$ | Understand | CO2 | AITB03.05 |
| Part - B (Long Answer Questions) |  |  |  |  |
| 1 | Convert Regular Expression 01* + 1 to Finite Automata. | Understand | CO2 | AITB03.05 |
| 2 | Construct Right linear, Left linear Regular Grammars for 01* +1 . | Apply | CO2 | AITB03.05 |
| 3 | Demonstrate Regular expression? Simplify the following Regular Expression <br> i) $\varepsilon+1^{*}(011)^{*}\left(1^{*}(011)^{*}\right)^{*}=(1+011)^{*}$ <br> ii) $(0+11 * 0)+(0+11 * 0)(10+10 * 1) *(10+10 * 1) *=1 * 0(10+10 * 1) *$ | Understand | CO2 |  |
| 4 | Construct Regular grammar for the given Finite Automata. $(\mathrm{a}+\mathrm{b})^{*} \mathrm{~b}^{*}$. | Apply | CO 2 | AITB03.08 |
| 5 | Construct Regular grammar for the given Finite Automata 0 * $11(0+1)^{*}$ | Apply | CO2 | AITB03.08 |
| 6 | Demonstrate Regular expression, Regular set and Finite Automata distinguish those with example representations. | Understand | CO2 | AITB03.08 |
| 7 | Construct the Finite Automata(NFA- $\varepsilon$ ) for given regular expression $(0+1) * 00(0+1)$ * | Apply | CO2 | AITB03.07 |
| 8 | Convert Regular Expression (b+aa)*a* to Finite Automata(NFA- $\varepsilon$ ). | Understand | CO2 | AITB03.05 |
| 9 | State Pumping Lemma for Regular Languages with a suitable example. | Remember | CO2 | AITB03.07 |
| 10 | Convert given Regular expression (a* ${ }^{*}$ ) * to FA(NFA- $\mathcal{L}$ ). | Understand | CO2 | AITB03.08 |
| 11 | Convert the following automata into Regular expression $\mathrm{M}=\left(\left\{\mathrm{q}_{1}, \mathrm{q}_{2}, \mathrm{q}_{3}\right\},\{0,1\}, \delta, \mathrm{q}_{1},\left\{\mathrm{q}_{1}\right\}\right)$ where $\delta$ is given by $\left[\delta\left(q_{1}, 0\right)=\left\{q_{1}\right\}, \delta\left(q_{1}, 1\right)=\left\{q_{2}\right\}\right]$ $\left[\delta\left(q_{2}, 0\right)=\left\{q_{3}\right\}, \delta\left(q_{2}, 1\right)=\left\{q_{2}\right\}\right]$ $\left[\delta\left(q_{3}, 0\right)=\left\{q_{1}\right\}, \delta\left(q_{3}, 1\right)=\left\{q_{2}\right\}\right]$ | Remember | CO2 | AITB03.08 |
| 12 | Describe Pumping lemma. Prove that the language $\mathrm{L}=\left\{\mathrm{yy} / \mathrm{y}\right.$ belongs $\left.\{0,1\}^{*}\right\}$ is not regular. | Understand | CO2 | AITB03.08 |
| 13 | Demonstrate Regular grammar? Explain the types of regular grammar with examples. | Remember | CO2 | AITB03.07 |
| 14 | IIustrate the steps for conversion of regular grammar to finite automata? Construct the FA for the following grammar $\begin{aligned} & \mathrm{S} \rightarrow \mathrm{aS} / \mathrm{bA} / \mathrm{b} \\ & \mathrm{~A} \rightarrow \mathrm{aA} / \mathrm{bS} / \mathrm{a} \end{aligned}$ | Understand | CO2 | AITB03.05 |
| 15 | Convert the given Regular Expression 1(11) * to FA and convert it in to NFA. | Remember | CO2 | AITB03.07 |
| 16 | Prove that the following languages is not regular <br> i) $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}} / \mathrm{n}>=1\right\}$ <br> ii) $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{P}} / \mathrm{p}\right.$ is prime $\}$ | Evaluate | CO2 | AITB03.08 |
| 17 | Convert the following regular expression to Regular grammar $(0+1) * 00(0+1) *$ | Remember | CO2 | AITB03.08 |
| 18 | Construct the Left Linear Grammar for the strings start with a over an alphabet $\sum=\{\mathrm{a}, \mathrm{b}\}$. | Apply | CO2 | AITB03.08 |
| 19 | Illustrate the steps for conversion from Finite Automata to Regular Expression with example? | Understand | CO2 | AITB03.07 |
| 20 | Describe Pumping lemma. Prove that the language $\mathrm{L}=\left\{\mathrm{yy} / \mathrm{y}\right.$ belongs $\left.\{0,1\}^{*}\right\}$ is not regular. | Understand | CO2 | AITB03.05 |


| 1 | Convert Regular Expression (11+0)*(00+1)* to Finite Automata. | Remember | CO2 | AITB03.08 |
| :---: | :---: | :---: | :---: | :---: |
| 2 | Construct Regular Grammar for the following Expressions <br> i) $a(a+b)^{*}$ <br> ii) $(a a+b b)$ | Understand | CO2 | AITB03.08 |
| 3 | DemonstratePumping Lemma for Regular Languages. Prove that the language $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} / \mathrm{n}\right.$ is $\left.\mathrm{an}^{5}\right\}$ is not regular | Understand | CO2 | AITB03.08 |
| 4 | Construct the DFA Transition diagram for equivalent Regular expression $(a b+a) *(a a+b)$ | Apply | CO 2 | AITB03.05 |
| 5 | Prove that following languages are not regular $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{m}} \mid \mathrm{n}, \mathrm{m}\right.$ andn $\left.<\mathrm{m}\right\}$ $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mid \mathrm{n}\right.$ is a perfect square $\}$ | Evaluate | CO2 | AITB03.06 |
| 6 | Construct the equivalent DFA for following Regular Expression $(0+1) *(00+11)(0+1)^{*}$ and Find reduced DFA. | Apply | CO2 | AITB03.08 |
| 7 | Construct the NFA for following Regular expression (0+1)*(01+110). | Apply | CO 2 | AITB03.08 |
| 8 | Convert the following automata into Regular expression $\mathrm{M}=\left(\left\{\mathrm{q}_{1}, \mathrm{q}_{2}, \mathrm{q}_{3}\right\},\{0,1\}, \delta, \mathrm{q}_{1},\left\{\mathrm{q}_{2}, \mathrm{q}_{3}\right\}\right)$ where $\delta$ is given by $\left[\delta\left(q_{1}, 0\right)=\left\{q_{2}\right\}, \delta\left(q_{1}, 1\right)=\left\{q_{3}\right\}\right]$ <br> $\left[\delta\left(\mathrm{q}_{2}, 0\right)=\left\{\mathrm{q}_{1}\right\}, \delta\left(\mathrm{q}_{2}, 1\right)=\left\{\mathrm{q}_{3}\right\}\right]$ <br> $\left[\delta\left(q_{3}, 0\right)=\left\{q_{2}\right\}, \delta\left(q_{3}, 1\right)=\left\{q_{2}\right\}\right]$ | Remember | CO 2 | AITB03.05 |
| 9 | Construct the following language is not regular i) $L=\left\{a^{n} b a^{n} / n=0,1,2 \ldots \ldots.\right\}$ <br> ii) $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{2 \mathrm{n}} / \mathrm{n} \geq 0\right\}$ | Apply | CO2 | AITB03.06 |
| 10 | Construct the Left Linear Grammar and Right Linear grammar for the automata in which strings end with 101 over an alphabet $\sum=\{0,1\}$. | Apply | CO2 | AITB03.08 |
| 11 | Construct a regular expression for the following Finite Automata | Apply | CO2 | AITB03.06 |
| MUUULE -III |  |  |  |  |
| CONTEXT FREE GRAMMARS |  |  |  |  |
| Part - A (Short Answer Questions) |  |  |  |  |
| S. No | Questions | Blooms <br> Taxonomy <br> Level | COs | Course Learning Outcomes |
| 1 | Develop a context free grammar( CFG). | Apply | CO3 | AITB03.09 |
| 2 | Construct the parse tree with example. | Apply | CO3 | AITB03.09 |
| 3 | Compare and contrast the Rightmost derivation with Left most derivation with example. | Understand | CO3 | AITB03.10 |
| 4 | Demonstrate a short notes about leftmost derivation with example. | Understand | CO3 | AITB03.10 |
| 5 | Express any two applications of Context Free Grammar. | Understand | CO3 | AITB03.11 |
| 6 | Narrate the left sentential form? | Understand | CO3 | AITB03.12 |
| 7 | Express the different ways to derive a string from a CFG. | Understand | CO3 | AITB03.09 |
| 8 | Demonstrate the language generated by CFG or G? | Understand | CO3 | AITB03.12 |
| 9 | Demonstrate the concept of parse tree? | Understand | CO3 | AITB03.11 |
| 10 | Demonstrate the concept of subtree. | Understand | CO3 | AITB03.10 |
| 11 | If S->aSb \| aAb, A->bAa, A->ba. Find out the CFL | Understand | CO3 | AITB03.10 |
| 12 | Demonstrate the usage of normalization? | Understand | CO3 | AITB03.09 |
| 13 | Analyze the ambiguous grammar with example? | Analyze | CO3 | AITB03.09 |
| 14 | Evaluate the following grammar into regular grammar that generates the same language S-> AB | Evaluate | CO3 | AITB03.09 |


|  | $\begin{array}{\|l} \hline \text { A->aAa\|bAb\|a\|b } \\ \text { B-> Ab\|Bb\| } \\ \hline \end{array}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 15 | Simplify the CFG to reduce UNIT production. | Analyze | CO3 | AITB03.11 |
| 16 | Express the elimination of useless symbols in productions. | Understand | CO3 | AITB03.12 |
| 17 | Solve the given grammar to get the minimized CFG - S $\rightarrow$ a S/A, A ->a / B | Apply | CO3 | AITB03.11 |
| 18 | Express the ambiguity concept in CFG with an example | Understand | CO3 | AITB03.10 |
| 19 | Describe the is the use of CNF and GNF. | Understand | CO3 | AITB03.11 |
| 20 | Evaluate the minimization of CFG - $\mathrm{S} \rightarrow \mathrm{aS} 1 \mathrm{bS} 1 \rightarrow \mathrm{aS} 1 \mathrm{~b} / \varepsilon$. | Evaluate | CO3 | AITB03.11 |
| 21 | Express the minimized CFG - $\mathrm{S} \rightarrow \mathrm{A}, \mathrm{A} \rightarrow \mathrm{aA} / \varepsilon$. | Understand | CO3 | AITB03.12 |
| 22 | Construct the minimized CFG - $\mathrm{S} \rightarrow \mathrm{AB} / \mathrm{a}, \mathrm{A} \rightarrow \mathrm{a}$. $\mathrm{B} \rightarrow \mathrm{b}$ | Apply | CO3 | AITB03.12 |
| 23 | Covert to to the minimized CFG - $\rightarrow$ aS/A/C A $\rightarrow \mathrm{aB} \rightarrow \mathrm{aaC} \rightarrow \mathrm{aCb}$. | Apply | CO3 | AITB03.12 |
| 24 | Convert the given grammar to CNF - $\mathrm{S} \rightarrow \mathrm{aAbB} \mathrm{A} \rightarrow \mathrm{aA} / \mathrm{a} \mathrm{B} \rightarrow \mathrm{bB} / \mathrm{a}$. | Understand | CO3 | AITB03.10 |
| Part - B (Long Answer Questions) |  |  |  |  |
| 1 | Construct Leftmost Derivation., Rightmost Derivation, Derivation Tree for the following grammar with respect to the string aaabbabbba. $\mathrm{S} \rightarrow \mathrm{aB}\|\mathrm{bA} \mathrm{~A} \rightarrow \mathrm{aS}\| \mathrm{bAA}\|\mathrm{aB} \rightarrow \mathrm{bS}\| \mathrm{aBB} \mid \mathrm{b}$ | Apply | CO3 | AITB03.12 |
| 2 | Design a CFG for the languages $\mathrm{L}=\left\{\mathrm{a}^{\left.\mathrm{i} \mathrm{b}^{\mathrm{j}} \mid \mathrm{i}<=2 \mathrm{j}\right\}}\right.$ | Create | CO3 | AITB03.12 |
| 3 | Construct leftmost and rightmost derivations for the strings, if the language is given as $\mathrm{S} \rightarrow \mathrm{AS} \mid \varepsilon$ <br> $\mathrm{A} \rightarrow$ aa\|ab|ba|bb <br> Strings: <br> a) aabbba <br> b) baabab <br> c) aaabbb | Apply | CO3 | AITB03.12 |
| 4 | Create the minimization of $\mathrm{CFG}-\mathrm{S} \rightarrow \mathrm{AbA} \mathrm{A} \rightarrow \mathrm{Aa} / \varepsilon$. | Create | CO3 | AITB03.12 |
| 5 | Calculate the minimization of CFG - $\mathrm{S} \rightarrow \mathrm{aSaS} \rightarrow \mathrm{bSb} \mathrm{S} \rightarrow \mathrm{a} / \mathrm{b} / \varepsilon$. | Understand | CO3 | AITB03.10 |
| 6 | Construct the minimization of CFG - $\mathrm{S} \rightarrow \mathrm{A} 0 / \mathrm{B} \mathrm{A} \rightarrow 0 / 12 / \mathrm{B}$ | Apply | CO3 | AITB03.11 |
| 7 | Convert the grammar to CNF - $\mathrm{S} \rightarrow \mathrm{aSa} / \mathrm{aas} \mathrm{S} \rightarrow \mathrm{bSb} / \mathrm{bb} \mathrm{S} \rightarrow \mathrm{a} / \mathrm{b}$. | Understand | CO3 | AITB03.12 |


| 8 | Write short notes on Chomsky Normal Form and Greibach Normal Form. | Understand | CO3 | AITB03.10 |
| :---: | :---: | :---: | :---: | :---: |
| 9 | What is Normalization of CFG? What is the use of Normalization? Explain different types of normal forms. | Remember | CO3 | AITB03.10 |
| 8 | Illustrate the construction of Greibach normal form with an example. | Understand | CO3 | AITB03.10 |
| 9 | Prove that the following CFG ambiguous. $\mathrm{S} \rightarrow \mathrm{iCtS}\|\mathrm{iCtSeS}\| \mathrm{a}, \mathrm{C} \rightarrow \mathrm{b}$. | Evaluate | CO3 | AITB03.12 |
| 10 | Demonstrate the Pumping lemma for Context Free Languages concept with example $\left\{a^{n} b^{n} c^{n}\right.$ where $\left.n>=0\right\}$. | Understand | CO3 | AITB03.12 |
| 11 | Illustrate the simplified CFG productions in $\mathrm{S} \rightarrow \mathrm{a}$ S1b S1 $\rightarrow$ a S1b/ $€$ | Understand | CO3 | AITB03.11 |
| 12 | Convert the following CFG into GNF. $\mathrm{S} \rightarrow \mathrm{AA} / \mathrm{a}, \mathrm{A} \rightarrow \mathrm{SS} / \mathrm{b}$ | Remember | CO3 | AITB03.11 |
| 13 | Describe unit production? Explain the procedure to eliminate unit production. | Understand | CO3 | AITB03.12 |
| 14 | Illustrate the procedure to eliminate $\epsilon$-productions in grammar. | Understand | CO3 | AITB03.09 |
| 15 | ```Convert the following grammar into GNF G=(\{A1,A2,A3\}, \(\{\mathrm{a}, \mathrm{b}\}, \mathrm{P}, \mathrm{A})\) \(\mathrm{A} 1 \rightarrow \mathrm{~A} 2 \mathrm{~A} 3\) \(\mathrm{A} 2 \rightarrow \mathrm{~A} 3 \mathrm{~A} 1 / \mathrm{b}\) \(\mathrm{A} 3 \rightarrow \mathrm{~A} 1 \mathrm{~A} 2 / \mathrm{a}\)``` | Understand | CO3 | AITB03.09 |
| 16 | ```Express simplified CFG productions from the following grammar \(\mathrm{A} \rightarrow \mathrm{aBb} / \mathrm{bBa}\) \(\mathrm{B} \rightarrow \mathrm{aB} / \mathrm{bB} / \epsilon\)``` | Understand | CO3 | AITB03.10 |
| 17. | Convert the following grammar into GNF $\mathrm{S} \rightarrow \mathrm{ABA} / \mathrm{AB} / \mathrm{BA} / \mathrm{AA} / \mathrm{B}$ $A \rightarrow a A / a, B \rightarrow b B / b$ | Understand | CO3 | AITB03.11 |
| 18 | Express the minimized CFG for the following grammar $\mathrm{S} \rightarrow \mathrm{ABCa} \mid \mathrm{bD}$ $\mathrm{A} \rightarrow \mathrm{BC}\|\mathrm{bB} \rightarrow \mathrm{b}\| \varepsilon$ $\mathrm{C} \rightarrow \mathrm{\oplus} \mid \varepsilon, \mathrm{D} \rightarrow \mathrm{~d}$ | Understand | CO3 | AITB03.12 |
| 19 | Covert the CFG to Greiback Normal form by taking an example | Understand | CO3 | AITB03.12 |
| 20 | $\begin{aligned} & \text { Design the grammar G given by } \\ & \text { S->aAa } \\ & \text { A->Sb\| bcc\|DaA } \\ & \text { C->abb\| DD } \\ & \text { E->ac } \\ & \text { D->aDa } \end{aligned}$ into an equivalent grammar by removing useless symbols and useless productions from it | Create | CO3 | AITB03.12 |
| Part - C (Problem Solving and Critical Thinking Questions) |  |  |  |  |
| 1 | Design a grammar for valid expressions over operator - and /. The arguments of expressions are valid identifiers over symbols a,b, 0 and Derive Left Most Derivation and Right Most Derivation for string W= (a11-b0) / (b00-a01). Draw parse tree for Left Most Derivation. | Understand | CO3 | AITB03.11 |
| 2 | Convert the following grammar into GNF $\mathrm{A} 1 \rightarrow \mathrm{~A} 2 \quad \mathrm{~A} 3 \mathrm{~A} 2 \rightarrow \mathrm{~A} 3 \mathrm{~A} 1 / \mathrm{b} A 3 \rightarrow \mathrm{~A} 1 \mathrm{~A} 2 / \mathrm{a}$ | Apply | CO3 | AITB03.12 |
| 3 | Evaluate the following grammar : $\mathrm{S} \rightarrow \mathrm{ABC} \mid \mathrm{BbB}$ $\mathrm{A} \rightarrow \mathrm{aA}\|\mathrm{BaC}\| \mathrm{aaa} \mathrm{~B} \rightarrow \mathrm{bBb}\|\mathrm{a}\| \mathrm{DC} \rightarrow \mathrm{CA} \mid \mathrm{AC}$ <br> $\mathrm{D} \rightarrow \varepsilon$ <br> Eliminate $\varepsilon$-productions. <br> Eliminate any unit productions in the resulting grammar. Eliminate any useless symbols in the resulting grammar. Convert the resulting grammar into Chomsky Normal Form | Evaluate | CO3 | AITB03.11 |
| 4 | Design a grammar for valid expressions over operator - and /. The Arguments of expressions are valid identifiers over symbols a,b, 0 and 1 . Derive Left Most Derivation and Right Most Derivation for string $\mathrm{W}=(\mathrm{a} 11-\mathrm{b} 0) /(\mathrm{b} 00-\mathrm{a} 01)$. Draw parse tree for Left Most Derivation. | Create | CO3 | AITB03.11 |
| 5 | Narrate the CFG G, find a CFG G" in Chomsky Normal form generating <br> L(G) - $\{\Lambda\}$ $\begin{aligned} & \mathrm{S} \rightarrow \mathrm{AaA}\|\mathrm{CA}\| \mathrm{BaB} \\ & \mathrm{~A} \rightarrow \mathrm{aaBa}\|\mathrm{CDA}\| \mathrm{aa} \mid \mathrm{DC} \\ & \mathrm{~B} \rightarrow \mathrm{bB}\|\mathrm{bAB}\| \mathrm{bb} \mid \mathrm{aS} \\ & \mathrm{C} \rightarrow C a\|b C\| D \\ & \hline \end{aligned}$ | Understand | CO3 | AITB03.11 |


|  | $\mathrm{D} \rightarrow b D \mid \Lambda$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 6 | Convert the following grammar into GNF $\mathrm{A} 1 \rightarrow \mathrm{~A} 2 \mathrm{~A} 3 \mathrm{~A} 2 \rightarrow \mathrm{~A} 3 \mathrm{~A} 1 / \mathrm{b}$ A3 $\rightarrow \mathrm{A} 1 \mathrm{~A} 2 / \mathrm{a}$ | Understand | CO3 | AITB03.11 |
| 7 | Describe CFG and Design a CFG for the following language. $\mathrm{L}=\left\{0^{\mathrm{i}} 1^{\mathrm{j}} 0^{\mathrm{k}} \mid j>i+k\right\}$ | Understand | CO3 | AITB03.11 |
| 8 | Explore the following grammar : <br> $S->A B C \mid B b B$ <br> A-> aA \|BaC|aaa <br> B-> bBb\|a|D <br> C->CA\|AC <br> D-> $\varepsilon$ <br> Eliminate $\varepsilon$-productions. <br> Eliminate any unit productions in the resulting grammar. <br> Eliminate any useless symbols in the resulting grammar. <br> Convert the resulting grammar into Chomsky Normal Form | Understand | CO3 | AITB03.11 |
| 9 | Express formally that the language $\left\{a^{n} b^{m} \mid n \leq m \leq 2 n\right\}$ is not deterministically context-free. Unfortunately, when applying the operation pre from the chapter, we obtain $\left\{a^{n} b^{m} \mid n<m \leq 2 n\right\}$ which still is context-free. We need a different closure property of the deterministic context-free languages to tackle this problem | Understand | CO3 | AITB03.11 |
| 10 | Demonstrate formally that the language $\left\{\operatorname{wir}(w) \mid w \in\{a, b\}^{*}\right\}$ is not deterministically context | Understand | CO3 | AITB03.12 |
| 11 | Show that $\left\{a^{m} b^{n} c^{p} \mid m<n\right.$ or $\left.n<p\right\}$ is not deterministically context | Understand | CO3 | AITB03.13 |
| 12 | Simplify the context free grammar for the given CFG $\begin{aligned} & \text { S-> Ab\|Bb } \\ & \text { A->a a aS \|Baa } \\ & \text { B-> b\|bS \|aBB } \end{aligned}$ | Analyze | CO3 | AITB03.11 |
| 14 | Construct the CFG for the language $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{2 \mathrm{n}}\right.$ where $\left.\mathrm{n}>=1\right\}$ | Understand | CO3 | AITB03.12 |
| 15 | Construct a CFG to generate unequal number of acs and bes | Understand | CO3 | AITB03.13 |
| 16 | Design the context free grammars in the four tuble form.(V,T,P,S) for the given languages on $\Sigma=\{a, b\}$ <br> i) All strings having at least two a "s <br> ii)All possible strings not containing triple $b^{\text {co }} s$ | Create | CO3 | AITB03.11 |
| 17 | Obtain the the string "aabbabba"for left most derivation and rightmost derivation using a CFG given by $\begin{aligned} & \text { S->Ab\|Ba } \\ & \text { A->a\|aS\|Baa } \\ & \text { B->b\|bS \|aBB } \\ & \hline \end{aligned}$ | Understand | CO3 | AITB03.11 |
| 18 | Describe and distinguish regular grammar and context free grammar | Understand | CO3 | AITB03.12 |
| 19 | Develop a left linear grammar for the Language given below $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{~b}^{\mathrm{m}} \mid \mathrm{n}>=2, \mathrm{~m}>=3\right\}$ | Apply | CO3 | AITB03.13 |
| 20 | Develop the context free grammar with the minimal number of variables that generates the language for the below languages $\begin{aligned} & \mathrm{L}=\left\{\mathrm{w} \mid \mathrm{w}=\mathrm{W}^{\mathrm{R}}\right\}\left(w^{\mathrm{R}} \text { denotes the reverse of } \mathrm{w}\right) \\ & \mathrm{L}=\left\{\mathrm{w} \mid \mathrm{w}=!\mathrm{W}^{\mathrm{R}}\right\} \end{aligned}$ | Apply | CO3 | AITB03.11 |
| MODULE-IV |  |  |  |  |
| PUSH DOWN AUTMATA |  |  |  |  |
| Part - A (Short Answer Questions) |  |  |  |  |
| 1. | Distinguish between deterministic and nondeterministic PDA. | Analyze | CO4 | AITB03.13 |
| 2. | Demonstrate the concept of PDA. | Understand | CO4 | AITB03.14 |
| 3. | Demonstrate the concept of NPDA. | Understand | CO4 | AITB03.14 |
| 4. | Describe the language of DPDA. | Understand | CO4 | AITB03.14 |
| 5. | Convert the following PDA to CFG $\delta(q 0,0, z 0)=\{q 0, x z 0)$ | Understand | CO4 | AITB03.15 |


| 6. | Convert the following PDA to CFG $\delta(\mathrm{q} 0,0, \mathrm{x})=(\mathrm{q} 0, \mathrm{xx})$ | Understand | CO4 | AITB03.15 |
| :---: | :---: | :---: | :---: | :---: |
| 7. | Convert the following PDA to CFG $\delta(\mathrm{q} 0,1, \mathrm{x})=(\mathrm{q} 1, \epsilon)$ | Understand | CO4 | AITB03.15 |
| 8. | Convert the following PDA to CFG $\delta(\mathrm{q} 1,1, \mathrm{x})=(\mathrm{q} 1, \mathrm{\epsilon})$ | Understand | CO4 | AITB03.15 |
| 9. | List out the steps to convert CFG to PDA. | Remember | CO4 | AITB03.15 |
| 10 | Express the acceptance of PDF by final state. | Understand | CO4 | AITB03.14 |
| 11 | Express the acceptance of PDF by empty stack. | Understand | CO4 | AITB03.14 |
| 12 | Convert the following PDA to CFG $\delta(\mathrm{q} 0, \mathrm{~b}, \mathrm{z} 0)=\{\mathrm{q} 0, \mathrm{zz} 0)$ | Understand | CO4 | AITB03.14 |
| 13. | construct the following PDA to CFG $\delta(\mathrm{q} 0, \mathrm{~b}, \mathrm{z})=(\mathrm{q} 0, \mathrm{zz})$ | Understand | CO4 | AITB03.14 |
| 14. | Convert the following PDA to CFG $\delta(\mathrm{q} 0, \epsilon, \mathrm{z} 0)=(\mathrm{q} 0, \epsilon)$ | Understand | CO4 | AITB03.15 |
| 15 | Describe the PDA and design PDA for $L=\left\{x \in\{a, b\}^{*} \mid\right.$ $n a(x)>n b(x)\}$ | Understand | CO4 | AITB03.15 |
| Part - B (Long Answer Questions) |  |  |  |  |
| 1. | State the NPDA(Nondeterministic PDA) and DPDA(deterministic PDA) equivalent? Illustrate with an example. | Remember | CO4 | AITB03.16 |
| 2. | Construct the grammar for the following PDA. $\mathrm{M}=(\{\mathrm{q} 0$, $\mathrm{q} 1\},\{0,1\},\{\mathrm{X}, \mathrm{z} 0\}, \mathrm{\delta}, \mathrm{q} 0, \mathrm{Z} 0, \Phi)$ and where $\delta$ is given by $\begin{aligned} & \delta(\mathrm{q} 0,0, \mathrm{z} 0)=\{(\mathrm{q} 0, \mathrm{XZ} 0)\}, \delta(\mathrm{q} 0,0, \mathrm{X})=\{(\mathrm{q} 0, \mathrm{XX})\}, \delta(\mathrm{q} 0,1, \mathrm{X})=\{(\mathrm{q} 1, \varepsilon)\}, \\ & \delta(\mathrm{q} 1,1, \mathrm{X})=\{(\mathrm{q} 1, \varepsilon)\}, \delta(\mathrm{q} 1, \varepsilon, \mathrm{X})=\{(\mathrm{q} 1, \varepsilon)\}, \delta(\mathrm{q} 1, \varepsilon, \mathrm{Z} 0)=\{(\mathrm{q} 1, \varepsilon)\} . \end{aligned}$ | Remember | CO4 | AITB03.15 |
| 3. | Construct PDA for string of form $\mathrm{a}^{\mathrm{n}} \mathrm{b}^{2 \mathrm{n}}$ | Apply | CO4 | AITB03.14 |
| 4. | Express PDA mathematically. With a neat diagram explain the working of a Turing Machine | Understand | CO4 | AITB03.15 |
| 5. | Evaluate the PDA that accepts the language $\left\{\mathrm{a}^{\wedge} \mathrm{m} \mathrm{b}^{\wedge} \mathrm{n} / \mathrm{n}>\mathrm{m}\right\}$ | Evaluate | CO4 | AITB03.15 |
| 6. | Design a PDA for the following grammar $\mathrm{S} \rightarrow 0 \mathrm{~A}, \mathrm{~A} \rightarrow 0 \mathrm{AB} / 1, \mathrm{~B} \rightarrow 1$ | Create | CO4 | AITB03.16 |
| 7. | ```Convert the following PDA to CFG M=(\{q0,q1\},\{a,b\},\{z0,za\}, \(\bar{\delta}, \mathrm{q} 0, \mathrm{z} 0, \Phi) \delta\) is given by, \(\delta(\mathrm{q} 0, \mathrm{a}, \mathrm{z} 0)=(\mathrm{q} 0, \mathrm{zz}) \delta(\mathrm{q} 0, \mathrm{a}, \mathrm{z})=(\mathrm{q} 0, \mathrm{zz} 0)\) \(\delta(q 0, b, z)=(q 1, \epsilon)\) \(\delta(q 1, b, z)=(q 1, \epsilon)\) \(\delta(q 1, \epsilon, z 0)=(q 1, \epsilon)\)``` | Understand | CO4 | AITB03.14 |
| 8. | Demonstrate the PDA mathematically. Construct the PDA for the following language. $\mathrm{L}=\left\{\mathrm{w} / \mathrm{w}\right.$ of form $\left.\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}}\right\}$. | Understand | CO4 | AITB03.14 |
| 9 | Evaluate and express the following For the language $L=\left\{x \operatorname{cxr} / \mathrm{x} €\{\mathrm{a}, \mathrm{b}\}^{*}\right\}$ design a PDA (Push Down Automata) and trace it for string "bacab" | Understand | CO4 | AITB03.16 |
| 10 | Demonstrate the Pushdown automaton A is specified by $\mathrm{A}=(\{\mathrm{q} 0, \mathrm{q} 1\},\{\mathrm{a}, \mathrm{b}\},\{\mathrm{Z}, \mathrm{X}\}, \delta, \mathrm{qin}, \mathrm{Z}, \varnothing)$, where $\delta$ contains the following transitions: $(\mathrm{q} 0, \mathrm{a}, \mathrm{Z})>(\mathrm{q} 0, \lambda),(\mathrm{q} 0, \mathrm{a}, \mathrm{Z})>\rightarrow(\mathrm{q} 0, \mathrm{XZin})$, $(\mathrm{q} 0, \mathrm{a}, \mathrm{X}) \rightarrow(\mathrm{q} 0, \mathrm{XX}),(\mathrm{q} 0, \mathrm{~b}, \mathrm{X}) \rightarrow(\mathrm{q} 1, \lambda)$, $(\mathrm{q} 1, \mathrm{~b}, \mathrm{X}) ~>(\mathrm{q} 1, \lambda),(\mathrm{q} 1, \mathrm{a}, \mathrm{Z}) ~) \rightarrow(\mathrm{q} 0, \mathrm{Z})$. <br> Infer a (reduced) context-free grammar G for the empty stack language of A, i.e., $\mathrm{L}(\mathrm{G})=\operatorname{Le}(\mathrm{A})$. | Understand | CO4 | AITB03.14 |
| 11 | Construct PDA for the below grammar as shown below $\begin{aligned} & \text { S->aABB \|aAA } \\ & \text { A->aBB \|a } \\ & \text { B->bBB \| A } \end{aligned}$ <br> that accepts the language generated by given grammar | Apply | CO4 | AITB03.14 |
| 12 | Design a PDA for the below CFG which generates the palindrome accepted by $L(G)$ <br> S->aSal bSb\|a|b | Create | CO4 | AITB03.14 |
| 13 | Design a PDA and describe a context free grammar for the language $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{a}} \mathrm{b}^{\mathrm{c}} \mathrm{c}^{\mathrm{k}} ; \mathrm{i}<\mathrm{j}\right.$ or $\left.\mathrm{j}<\mathrm{k}\right\}$ | Create | CO4 | AITB03.16 |
| 14 | Covert the following context free grammar to push down automata $\begin{array}{\|l} \text { S->aA\|bB } \\ \text { A-> aB\|a } \\ \hline \end{array}$ | Understand | CO4 | AITB03.14 |


|  | B->b <br> Verify the string aab accepted by equivalent PDA |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 15 | Design DPDA for $\mathrm{L}=\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}}$ where $\mathrm{n}>=1$ | Create | CO4 | AITB03.14 |
| 16 | Construct PDA accepts PDA M for the language $\mathrm{L}=\left\{\mathrm{WWR} \mid \mathrm{W} \in\{\mathrm{a}, \mathrm{b}\}^{*}\right\}$ such that $\mathrm{L}=\mathrm{L}(\mathrm{M})$ | Apply | CO4 | AITB03.14 |
| 17 | Design PDA M for the language $\mathrm{L}=\left\{\mathrm{x} \in\{\mathrm{a}, \mathrm{b}\}^{*} \mid \mathrm{n}_{\mathrm{a}}(\mathrm{x})>\mathrm{n}_{\mathrm{b}}(\mathrm{x})\right\}$ | Create | CO4 | AITB03.16 |
| 18 | Prove that the below languages are deterministic context free languages? <br> a) $\mathrm{L} 1=\left\{0^{\mathrm{n}} 1^{\mathrm{m}} \mid \mathrm{n}=\mathrm{m}\right.$ and $\left.\mathrm{n}>=1\right\}$ <br> b) $\mathrm{L} 2=\left\{\mathrm{o}^{\mathrm{n}} 1^{\mathrm{m}} \mid \mathrm{n}=2 \mathrm{~m}\right.$ and $\left.\mathrm{n}>=1\right\}$ | Evaluate | CO4 | AITB03.14 |
| 19 | Narrate deterministic context free languages and deterministic push down automata | Understand | CO4 | AITB03.14 |
| 20 | Construct PDA that recognizes the language $\mathrm{L}=\left\{\mathrm{x}=\mathrm{x}^{\mathrm{R}}: \mathrm{x} \in\{\mathrm{a}, \mathrm{b}\}^{+}\right\}$ | Apply | CO4 | AITB03.14 |
| Part - C (Problem Solving and Critical Thinking Questions) |  |  |  |  |
| 1 | Construct PDA for equal number of x's and y's. eg: xyyxxy | Apply | CO4 | AITB03.14 |
| 2 | Construct NDPDA for $\mathrm{L}=\left\{\mathrm{W} \# \mathrm{~W}^{\mathrm{R}} / \mathrm{W} \in(\mathrm{X}+\mathrm{Y})^{*}\right\}$ | Apply | CO4 | AITB03.14 |
| 3 | $\begin{aligned} & \text { Convert the following PDA to CFG } \delta(\mathrm{q} 0,0, \mathrm{z} 0)=\{\mathrm{q} 0, \mathrm{xz} 0) \\ & \delta(\mathrm{q} 0,0, \mathrm{x})=(\mathrm{q} 0, \mathrm{xx}) \\ & \delta(\mathrm{q} 0,1, \mathrm{x})=(\mathrm{q} 1, \epsilon) \\ & \delta(\mathrm{q} 1,1, \mathrm{x})=(\mathrm{q} 1, \epsilon) \\ & \delta(\mathrm{q} 1, \epsilon, \mathrm{x})=(\mathrm{q} 1, \boldsymbol{\epsilon}) \\ & \delta(\mathrm{q} 1, \epsilon, \mathrm{z} 0)=(\mathrm{q} 1, \epsilon) \\ & \hline \end{aligned}$ | Understand | CO4 | AITB03.15 |
| 4 | Construct DPDA for $\mathrm{L}=\left\{\mathrm{W} \# \mathrm{~W}^{\mathrm{R}} / \mathrm{W} \in(\mathrm{X}+\mathrm{Y})^{*}\right\}$ | Apply | CO4 | AITB03.15 |
| 5 | Construct pushdown automata for the following languages. Acceptance either by empty stack or by finalstate. <br> (a) $\left\{a^{n} b^{m} a^{n} \mid m, n \in \mathrm{~N}\right\}$ <br> (b) $\left\{a^{n} b^{m} c^{m} \mid m, n \in N\right\}$ <br> (c) aibjck $\mid \mathrm{i}, \mathrm{j}, \mathrm{k} \in \mathrm{N}, \mathrm{i}>\mathrm{j}\}$ <br> (d) $\left\{a^{i} b^{j} c^{k} \mid i, j, k \in \mathrm{~N}, i+j=k\right\}$ <br> (e) $\left\{a^{i} b^{j} c^{k} \mid i, j, k \in \mathrm{~N}, i+k=j\right\}$ <br> (f) $\left\{a^{n} b^{m} \mid n \leq m \leq 2 n\right\}$ <br> (g)PAL $=\left\{w \in\{a, b\}^{*} \mid \operatorname{mir}(w)=w\right\}$ <br> (h) $\left\{w_{1} c w_{2} c \cdots c w_{k} c x \mid x, w_{1}, \ldots, w_{k} \in\{a, b\}^{*}\right.$, $k \in \mathrm{~N}, x=\operatorname{mir}\left(w_{j}\right)$ for some $\left.j\right\}$ <br> (i) $\left\{w \in\{a, b\}^{*} \mid \#_{a}(w)=\#_{b}(w)\right\}$, <br> $\#_{a}(w)$ represents the number of $a^{\prime \prime}$ s in $w$ <br> (j) $\left\{w \in\{a, b\} * \mid \#_{a}(w)=2 \cdot \#_{b}(w)\right\}$ | Apply | CO4 | AITB03.15 |
| 6 | Construct a PDA with final state acceptance for thelanguage $\mathrm{B}=\{\operatorname{bin}(i) \$ \operatorname{mir}(\operatorname{bin}(i+1)) \mid i \geq 0\} \subseteq\{0,1, \$\}^{*}$ <br> Here is $\operatorname{bin}(\mathrm{i}) \in\{0,1\} *$ the binary representation (without leading zero"s) of the number i. <br> $\operatorname{Eg} . \operatorname{bin}(11)=1011$ and $\operatorname{mir}(\operatorname{bin}(12))=0011$ | Apply | CO4 | AITB03.15 |
| 7 | Construct CFG corresponding to PDA whose transition mapping is as follows. $\begin{aligned} & \delta(\mathrm{S}, \mathrm{a}, \mathrm{X})=(\mathrm{s}, \mathrm{~A}, \mathrm{X}) \\ & \delta(\mathrm{S}, \mathrm{~b}, \mathrm{~A})=(\mathrm{s}, \mathrm{AA}) \\ & \delta(\mathrm{S}, \mathrm{a}, \mathrm{~A})=(\mathrm{s}, \mathrm{AA}) \end{aligned}$ | Apply | CO4 | AITB03.14 |


| 8 | Prove that given CFG with following productions $\mathrm{S}->\mathrm{aBc}$ $\mathrm{A}->\mathrm{abc}$ $\mathrm{B}->\mathrm{aAb}$ $\mathrm{C}->\mathrm{AB}$ $\mathrm{C}->\mathrm{c}$ constructs a PDA M such that the language generated by M and G are equivalent. | Evaluate | CO4 | AITB03.14 |
| :---: | :---: | :---: | :---: | :---: |
| 9 | $\begin{aligned} & \text { Design a PDA for the following grammar. } \\ & \text { S }->0 \mathrm{~A} \\ & \text { A->AB } \\ & B->1 \end{aligned}$ | Create | CO4 | AITB03.15 |
| 10 | Construct PDA for the following grammar $\begin{aligned} & \text { S->AA \|a } \\ & \text { A->SA \|b } \end{aligned}$ | Create | CO4 | AITB03.15 |
| MODULE-V |  |  |  |  |
| TURING MACHINE |  |  |  |  |
| Part - A (Short Answer Questions) |  |  |  |  |
| S. No | Questions | Blooms Taxonomy Level | COs | Course Learning Outcomes |
| 1 | Describe the Chomsky hierarchy of languages. | Understand | CO5 | AITB03.16 |
| 2 | Demonstrate Context sensitive language. | Understand | CO5 | AITB03.16 |
| 3 | Explore Turing Machine | Understand | CO5 | AITB03.16 |
| 4 | Express Type 0 grammars . | Understand | CO5 | AITB03.16 |
| 5 | Narrate the Type 1 grammars . | Understand | CO5 | AITB03.16 |
| 6 | Demonstrate Type 2 grammars . | Understand | CO5 | AITB03.16 |
| 7 | Describe Type 3 grammars . | Remember | CO5 | AITB03.16 |
| 8 | List out the types of grammars. | Remember | CO5 | AITB03.16 |
| 9 | Describe the moves in Turing Machine. | Understand | CO5 | AITB03.16 |
| 10 | Demonstrate an Instantaneous Description of a Turing Machine. | Remember | CO5 | AITB03.17 |
| 11 | Express the Language of Turing Machine. | Remember | CO5 | AITB03.17 |
| 12 | List out types of TMs. | Remember | CO5 | AITB03.18 |
| 13 | Distinguish the difference between PDA and TM | Understand | CO5 | AITB03.17 |
| 14 | Describe the multi head Turing Machine. | Understand | CO5 | AITB03.18 |
| 15 | obtain multi dimensional Turing Machine. | Remember | CO5 | AITB03.18 |
| 16 | Demonstrate multiple tapes Turing Machine. | Understand | CO5 | AITB03.18 |
| 17 | Describe the recursive languages. | Remember | CO5 | AITB03.17 |
| 18 | Construct recursively enumerable languages. | Apply | CO5 | AITB03.17 |
| 19 | Describe two way infinite Turing Machine. | Remember | CO5 | AITB03.18 |
| 20 | Demonstrate the non deterministic Turing Machine. | Understand | CO5 | AITB03.18 |
| 21 | Construct Turing Machine for 1's complement for binary numbers. | Understand | CO5 | AITB03.18 |
| 22 | Distinguish Recursive languages and Recursively enumerable languages. | Understand | CO5 | AITB03.19 |
| 23 | Demonstrate Church's Hypothesis. | Understand | CO5 | AITB03.20 |
| Part - B (Long Answer Questions) |  |  |  |  |
| 1 | Describe short notes on Context sensitive language and linear bounded automata. | Understand | CO5 | AITB03.16 |
| 2 | Classify briefly about Chomsky hierarchy of languages.. | Understand | CO5 | AITB03.16 |
| 3 | Describe a Turing Machine. With a neat diagram explain the working of a Turing Machine. | Understand | CO5 | AITB03.16 |
| 4 | Compare Turing Machine with other automata. | Understand | CO5 | AITB03.18 |
| 5 | Construct a Transition diagram for Turing Machine to accept the language $\mathrm{L}=\left\{\mathrm{w}^{\mathrm{w}} \mathrm{w}^{\mathrm{R}} \mid \mathrm{w} \in(\mathrm{a}+\mathrm{b}) *\right\}$ | Understand | CO5 | AITB03.17 |
| 6 | Express short notes on Recursive and Recursively Enumerable languages. | Understand | CO5 | AITB03.17 |


| 7 | Describe the properties of recursive and recursively enumerable <br> languages. | Understand | CO5 | AITB03.17 |
| :---: | :--- | :---: | :---: | :---: |
| 8 | Develop a Turing Machine to accept strings formed with 0 and <br> and having substring 000. | Understand | CO5 | AITB03.16 |
| 9 | Construct a Transition diagram for Turing Machine to accept the <br> language $\mathrm{L}=\left\{\mathrm{ww}^{\mathrm{R}} \mid \mathrm{w} \epsilon(\mathrm{a}+\mathrm{b}) *\right\}$ | Apply | CO5 | AITB03.16 |
| 10 | Design a Transition table for TM $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}} \mathrm{c}^{n} / \mathrm{n}>=1\right\}$ | Create | CO5 | AITB03.16 |
| 11 | Construct a Transition table for Turing Machine to accept the <br> following language. $\mathrm{L}=\left\{0^{\mathrm{n}} 1^{n} 0^{n} \mid \mathrm{n} \geq 1\right\}$ | Apply | CO5 | AITB03.16 |
| 12 | Construct a Turing Machine that accepts the language <br> $\mathrm{L}=\left\{1^{n} 2^{n} 3^{n} \mid n \geq 1\right\}$. Give the transition diagram for the Turing <br> Machine obtained and also show the moves made by the Turing <br> machine for the string 111222333. | Apply | CO5 | AITB03.16 |


| 13 | Enumerate Linear bounded automata and explain its model? | Understand | CO5 | AITB03.16 |
| :---: | :---: | :---: | :---: | :---: |
| 14 | Demonstrate the power and limitations of Turing machine. | Understand | CO5 | AITB03.18 |
| 15 | Construct Transition diagram for Turing Machine - L=\{a $\left.\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}} \mathrm{c}^{1} / \mathrm{n}=1\right\}$ | Apply | CO5 | AITB03.16 |
| 16 | Construct a Transition diagram for Turing Machine to implement addition of two unary numbers $(\mathrm{X}+\mathrm{Y})$. | Apply | CO5 | AITB03.16 |
| 17 | Construct a Linear Bounded automata for a language where $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}} / \mathrm{n}>=1\right\}$ | Apply | CO5 | AITB03.16 |
| 18 | Classify the types of Turing machines. | Understand | CO5 | AITB03.18 |
| 19. | Describe briefly about the following <br> a)Church's Hypothesis <br> b)Counter machine | Understand | CO5 | AITB03.16 |
| 20 | Construct Transition diagram for Turing Machine that accepts the language $\mathrm{L}=\left\{0^{\mathrm{n}} 1^{\mathrm{n}} \mid \mathrm{n} \geq 1\right\}$. Give the transition diagram for the Turing Machine obtained and also show the moves made by the Turing machine for the string 000111. | Apply | CO5 | AITB03.16 |
| Part - C (Problem Solving and Critical Thinking Questions) |  |  |  |  |
| 1 | Construct a Turing Machine that accepts the language $\mathrm{L}=\left\{\mathrm{a}^{2 \mathrm{n}} \mathrm{b}^{\mathrm{n}} \mid \mathrm{n} \geq 0\right\}$. Give the transition diagram for the Turing Machine obtained. | Apply | CO5 | AITB03.16 |
| 2 | Construct a Turing Machine that gives two's compliment for the given binary representation. | Apply | CO5 | AITB03.16 |
| 3 | Examine Type 3 and Type 2 grammars with example. | Analyze | CO5 | AITB03.15 |
| 4 | Extend the Type 1 and Type 0 grammars with example. | Understand | CO5 | AITB03.17 |
| 5 | Design a Turing Machine that accepts the set of all even palindromes over $\{0,1\}$ | Create | CO5 | AITB03.16 |
| 6 | Design Turing Machine for $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}} \mathrm{c}^{\mathrm{n}} \mid \mathrm{n} \geq 1\right\}$ | Create | CO5 | AITB03.16 |
| 7 | Construct Turing Machine to calculate GCD of two given numbers | Apply | CO5 | AITB03.15 |
| 8 | Compare and contrast the Finite state machine, PDA and Turing Machine | Understand | CO5 | AITB03.17 |
| 9 | Construct a Turing Machine to accept the following languages $\mathrm{L}=\left\{\mathrm{w}^{\mathrm{n}} \mathrm{x}^{\mathrm{n}} \mathrm{y}^{\mathrm{n}} \mathrm{n}^{\mathrm{n}} \mid \mathrm{n} \geq 1\right\}$ | Apply | CO5 | AITB03.16 |
| 10 | Design a Turing Machine that accepts the language denoted by regular expression (000)* | Create | CO5 | AITB03.16 |

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